

Improved Freezing In Vacuum

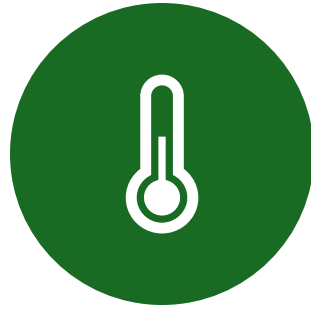
Melanie Plageman

Microsoft

Agenda



WHAT IS FREEZING AND
WHY DO WE NEED IT



WHEN DO WE FREEZE
NOW



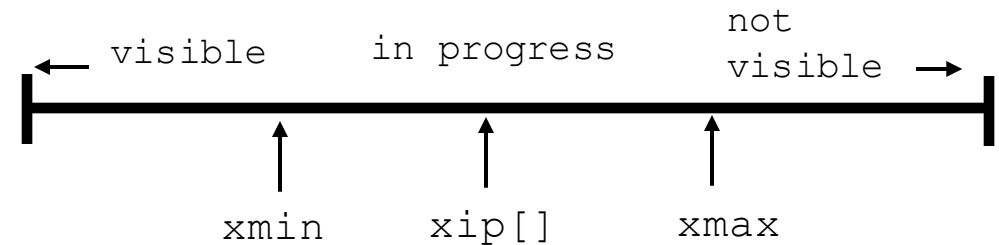
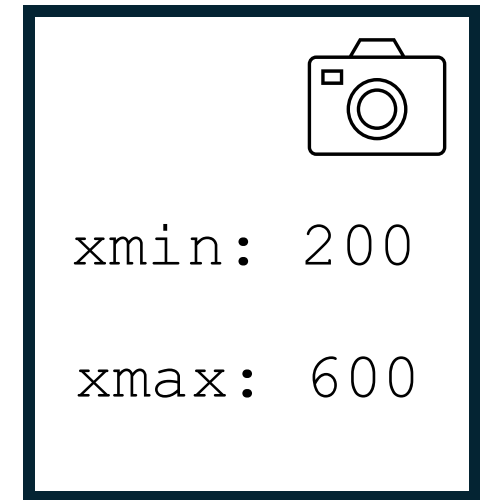
WHY IS IT HARD TO DO
BETTER



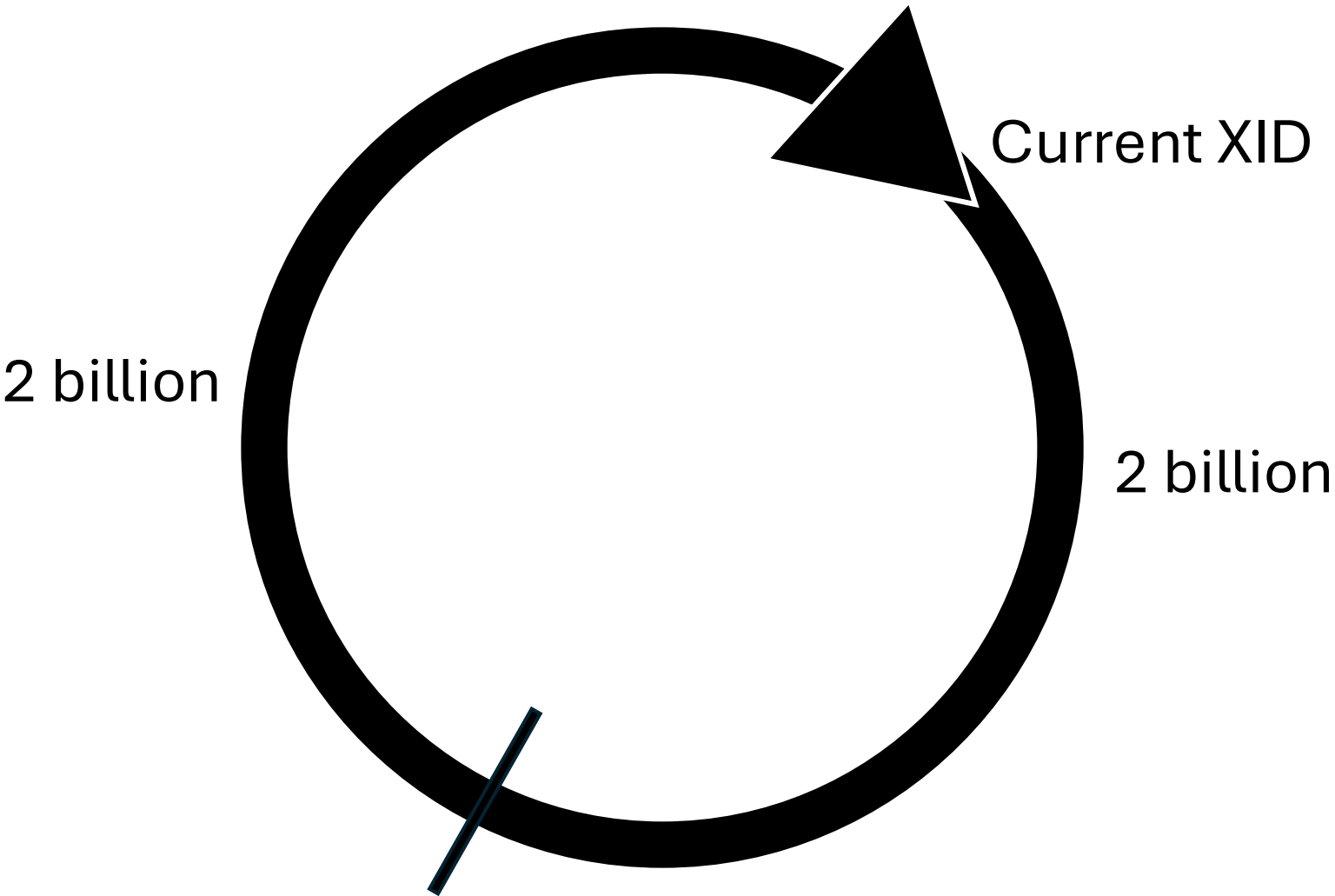
WHAT DID WE DO IN 18

Multi-Version Concurrency Control

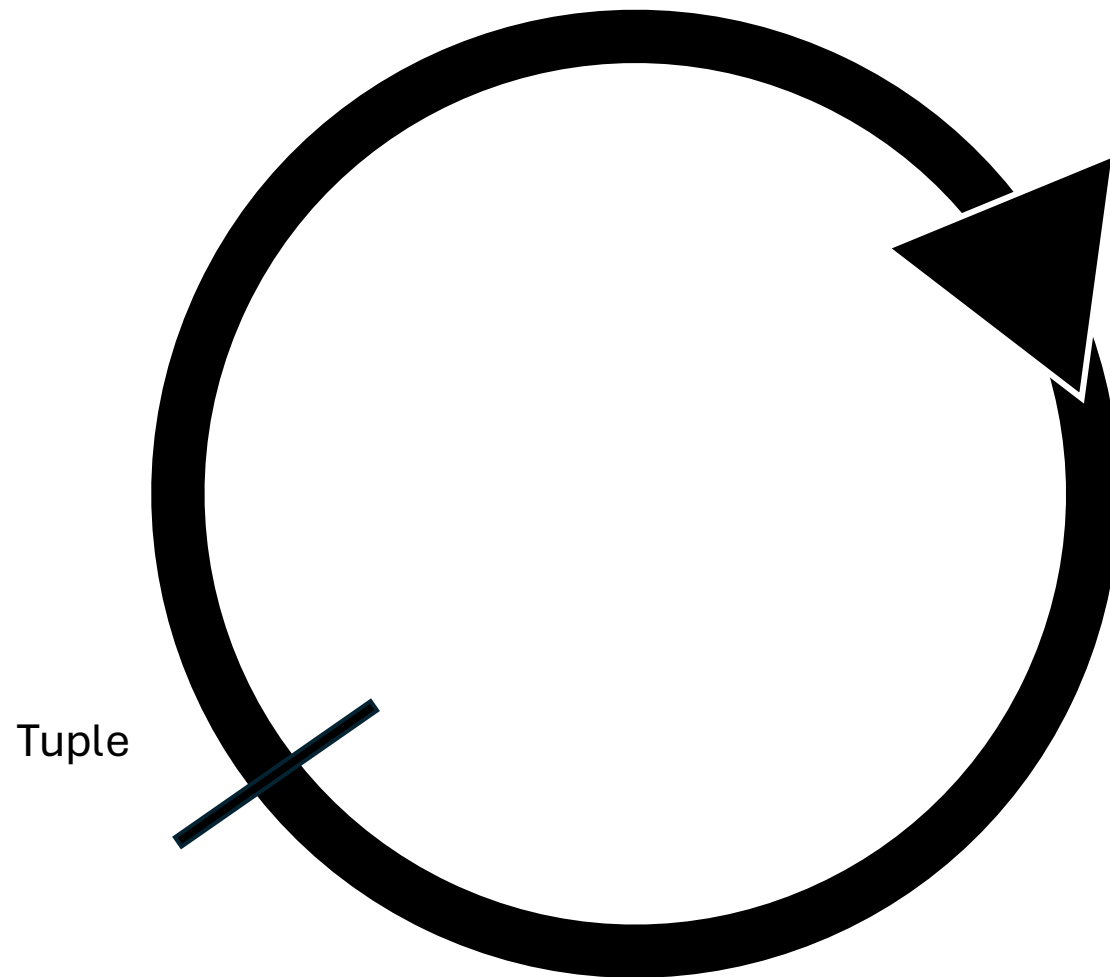
	xmin	xmax	status	value	
	20	X	R	lorem	Not visible, Live (running)
☠	22	100	C	ipsum	Visible, Dead
☠	50	700	C	dolor	Visible, Live; Not visible, Dead
	100	X	C	sit	Visible, Live
	300	X	C	amet	Check XIP



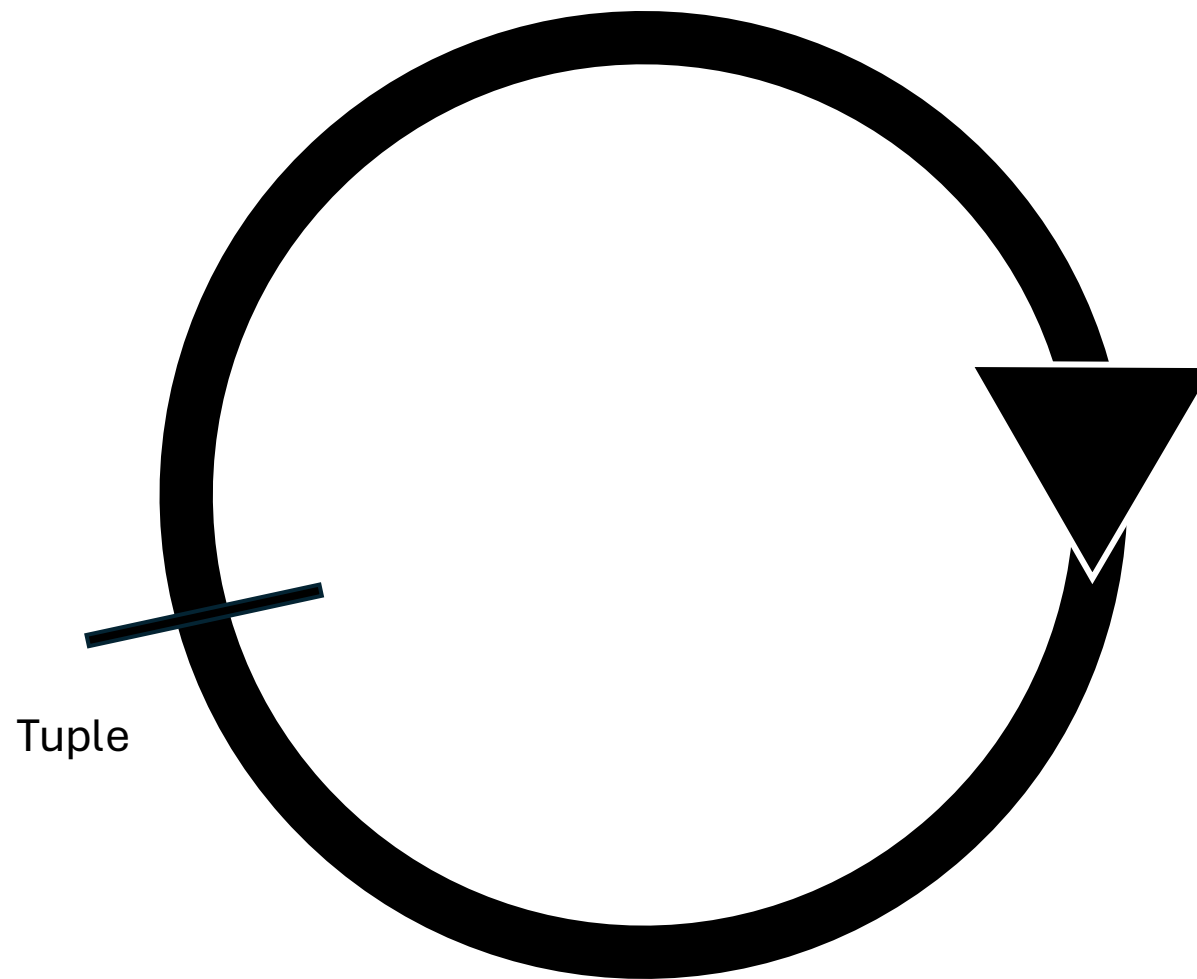
Wraparound



Wraparound



Wraparound



Freezing

	xmin	xmax	status	value	
	20	X	R	lorem	Not visible, Live (running)
	22	100	C	ipsum	Visible, Dead
	50	700	C	dolor	Visible, Live; Not visible, Dead
			C	sit	Visible, Live
	300	X	C	amet	Check XIP

Wraparound



Read-only database

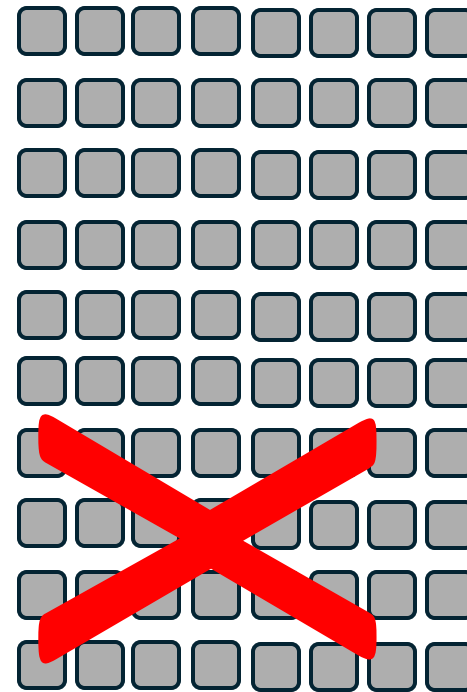
How do we know if our table is in danger?

Query XID: 1,900,000,000

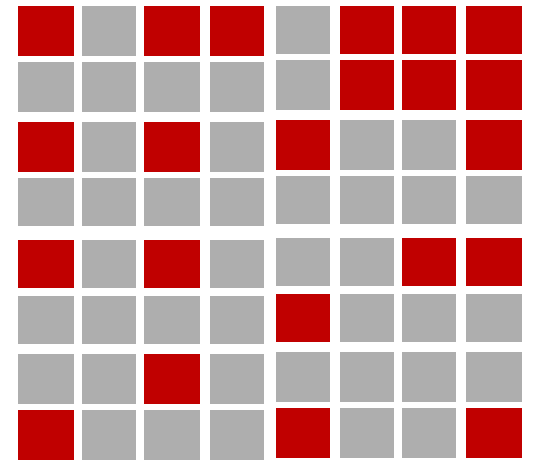


xmin	xmax	status	value
20	X	R	lorem
22	100	C	ipsum
50	700	C	dolor
100	X	C	sit
300	X	C	amet

`relfrozenxid == 10`



`relfrozenxid ==
1,000,000,000`



`relfrozenxid ==
1,200,000,000`

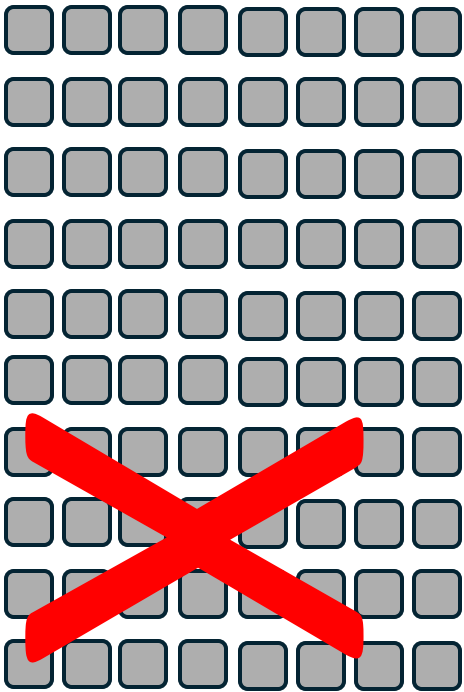
pg_class.relrozenxid

Query XID: 1,900,000,000

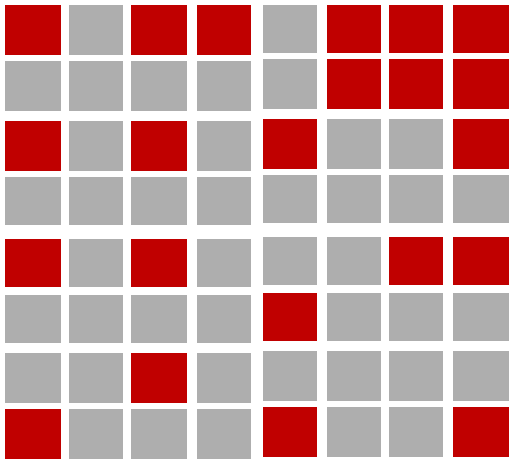


xmin	xmax	status	value
20	X	R	lorem
22	100	C	ipsum
50	700	C	dolor
100	X	C	sit
300	X	C	amet

relrozenxid == 10



relrozenxid ==
1,000,000,000



relrozenxid ==
1,200,000,000

Who should freeze and when?

Emits
WAL

Dirtyes
pages

Vacuum well-positioned to freeze



Encounters
tuples

Dirties
pages

Emits WAL



Freezing inconsistent with vacuum's mandate

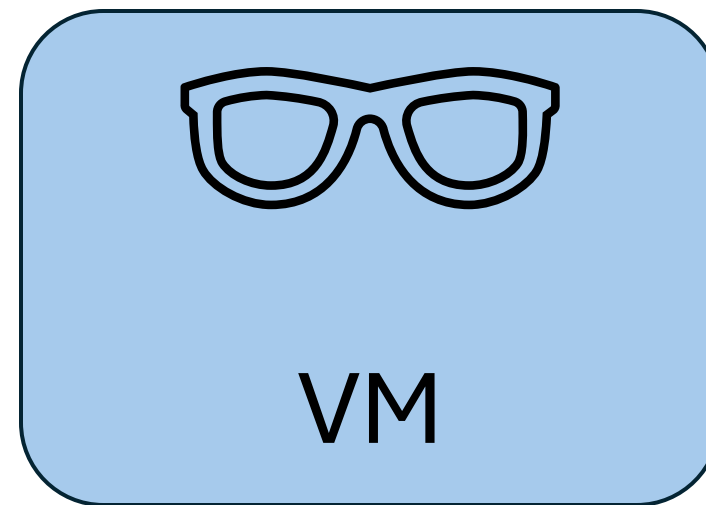
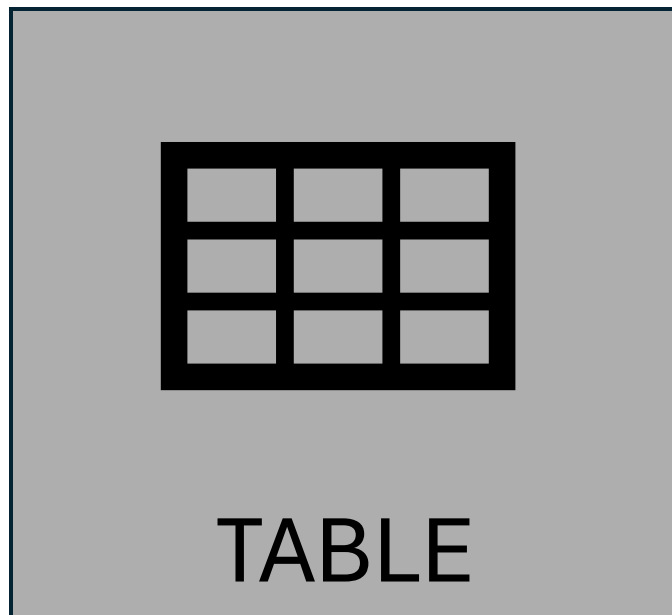


Triggered by
insert and
modification
thresholds

Reads modified
pages

Skips all-visible
pages

Visibility map



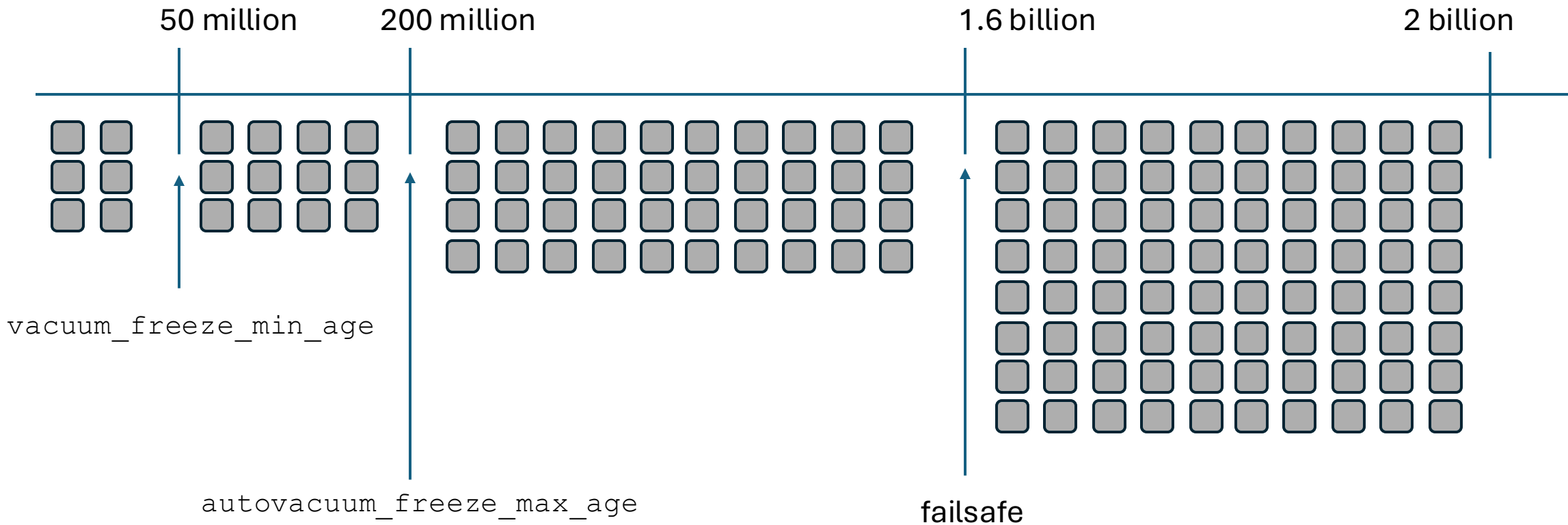
Forcing vacuum of all-visible pages

`relfrozenxid < autovacuum_freeze_max_age`

triggers anti-wraparound vacuum

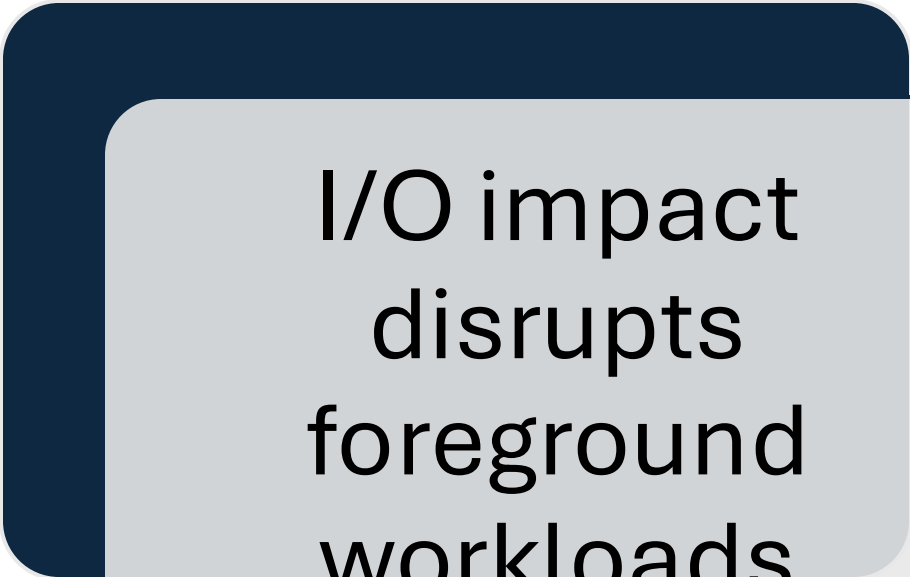
- usually aggressive (controlled by `vacuum_freeze_table_age`)
- scans all all-visible pages

Freezing Timeline

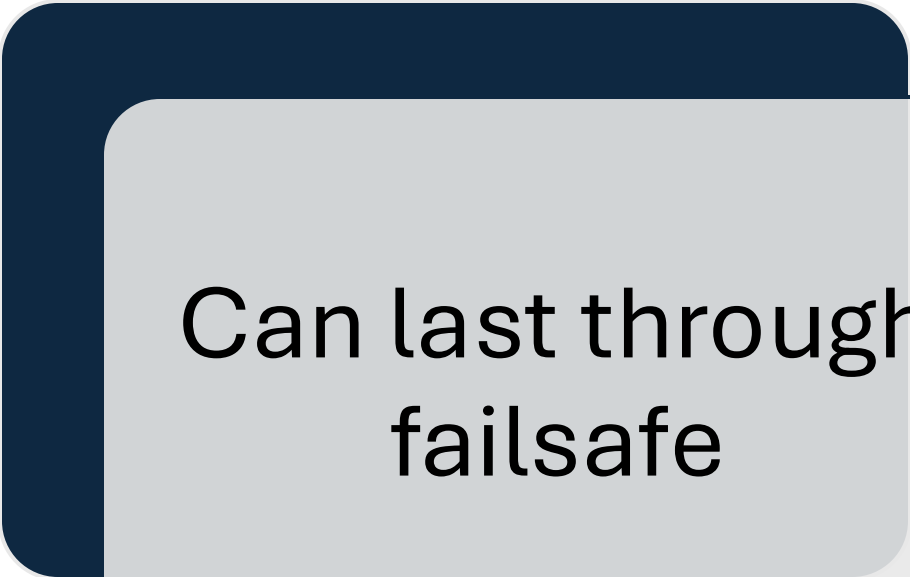




Aggressive Vacuum Overhead



I/O impact
disrupts
foreground
workloads



Can last through
failsafe

I/O Amplification Waiting to Freeze

Evict inserted
data



Write

SELECT * query



Read

Set
hints



C

Evict hinted
data



Write

Normal
vacuum



Read

Set page vis
hint and VM



PD_ALL_VISIBLE



Vacuum
strategy evict



Write

Aggressive
vacuum



Read

Freeze tuples
and update VM



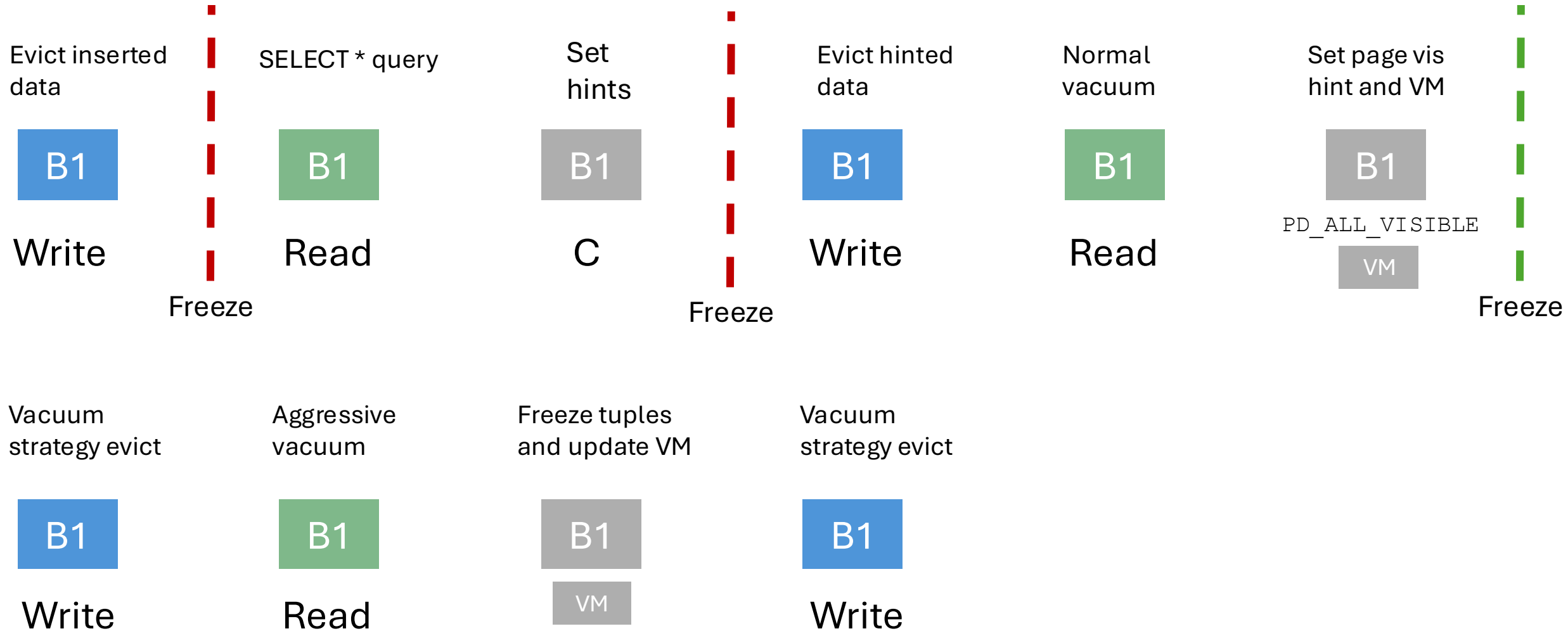
Vacuum
strategy evict



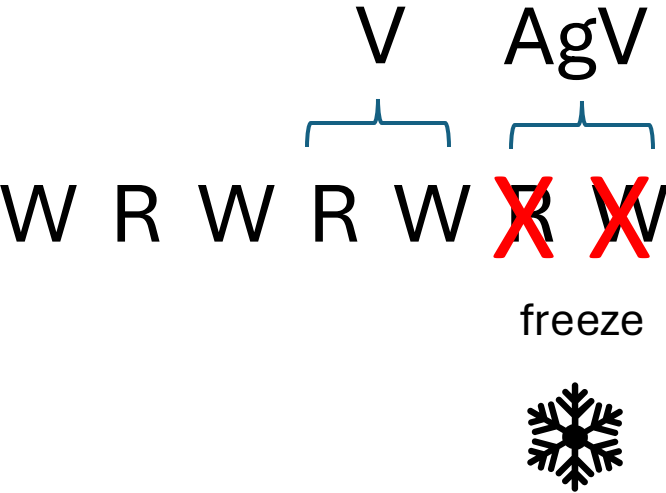
Write

When *should* we freeze?

Ideal Time To Freeze

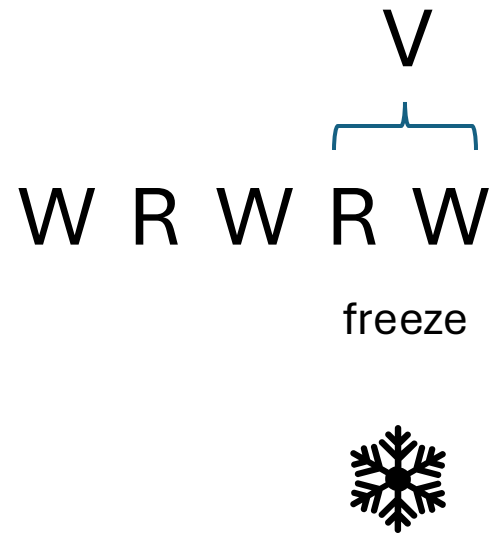


What can be eliminated

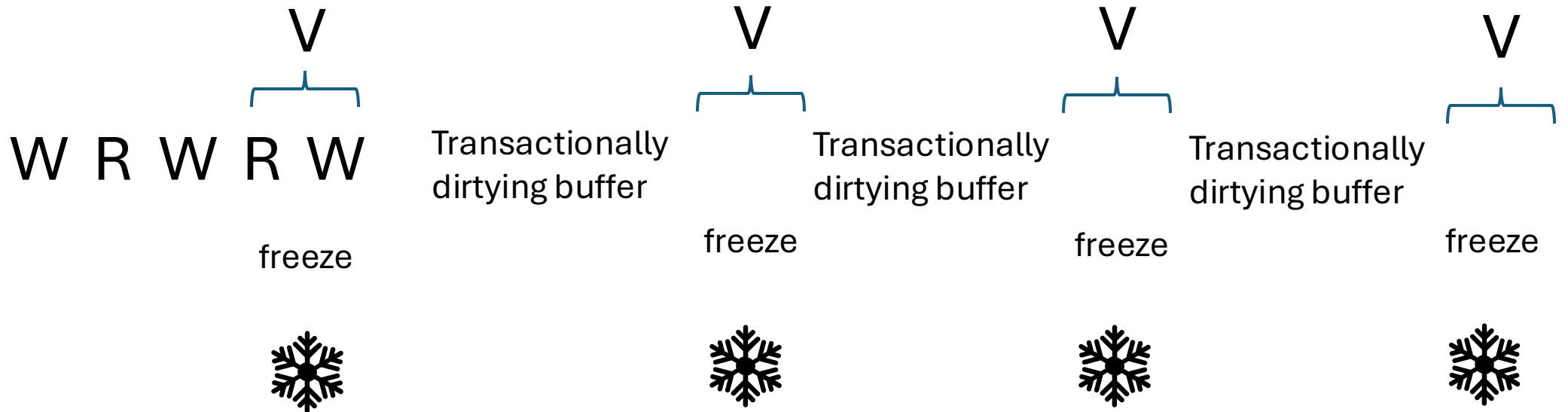


Preemptive Freeze Algorithm Possibilities

Freeze all visible tuples on every vacuum? (Ignore vacuum_freeze_min_age)

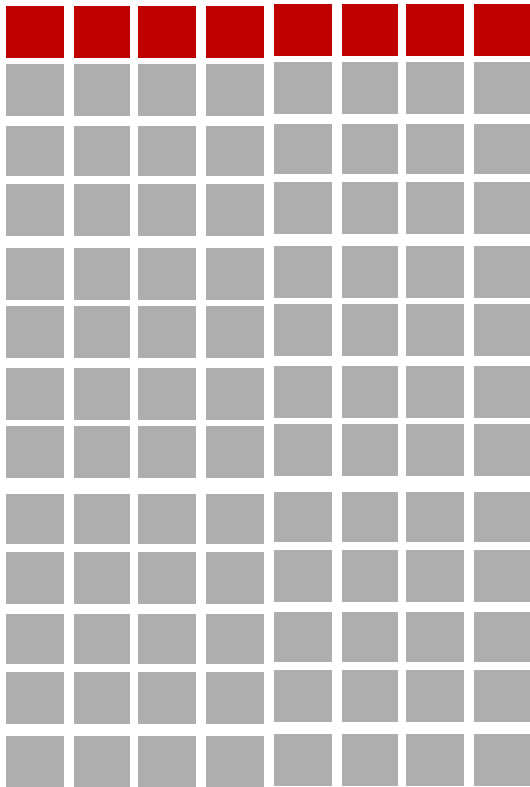


~~Freeze all visible tuples on every vacuum?~~

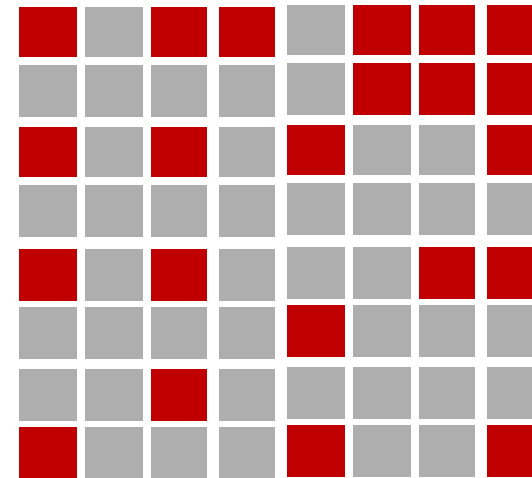


Diverse access patterns

Insert-only table



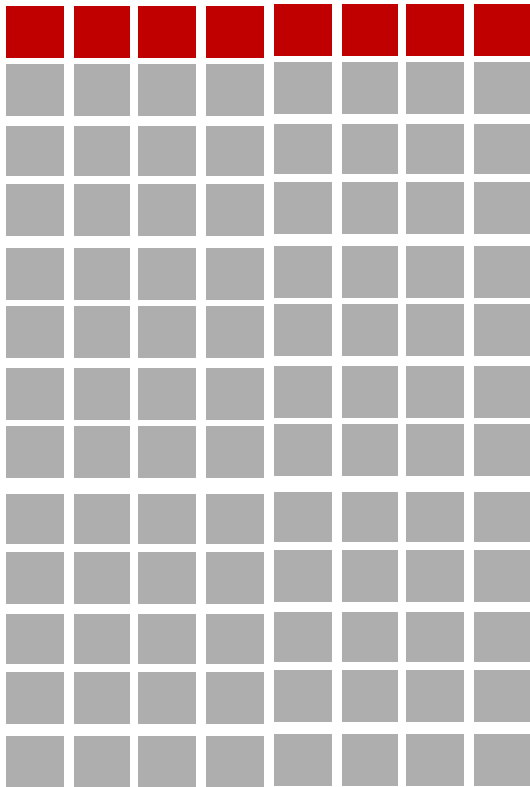
Hotly updated table



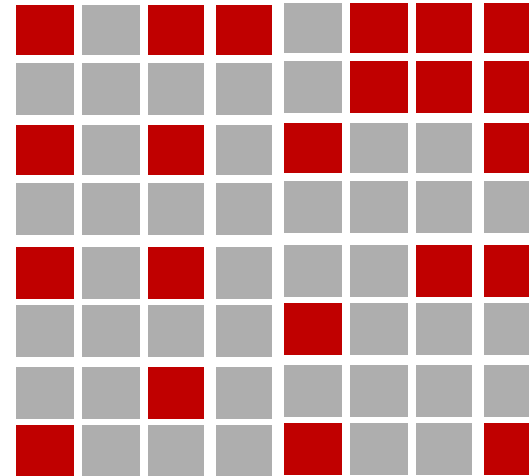
What about using # updates/deletes

```
pg_stat_all_tables.n_tup_upd, n_tup_del
```

Insert-only table



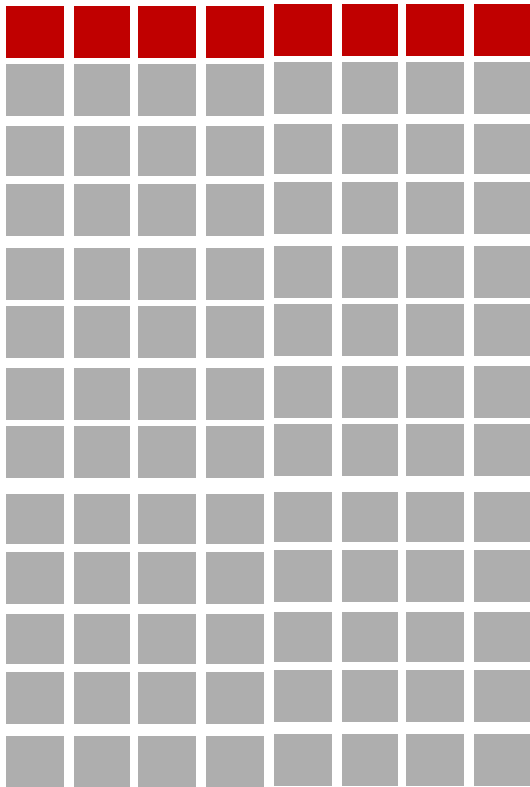
Hotly updated table



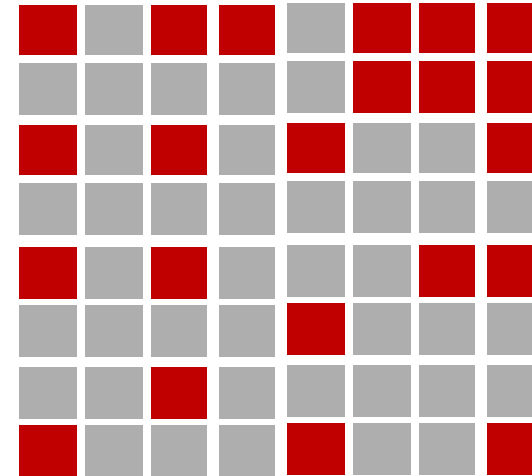
~~What about using # updates/deletes~~

```
pg_stat_all_tables.n_tup_upd, n_tup_del
```

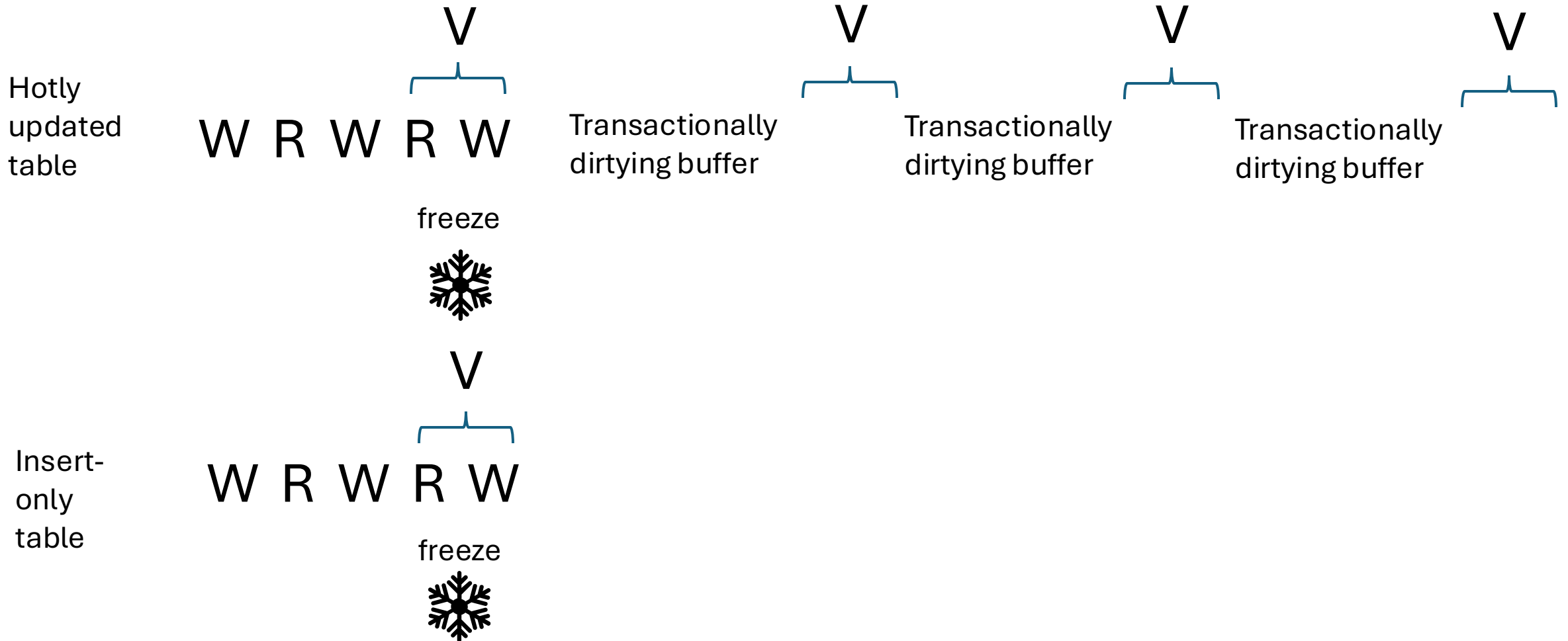
Insert-only table



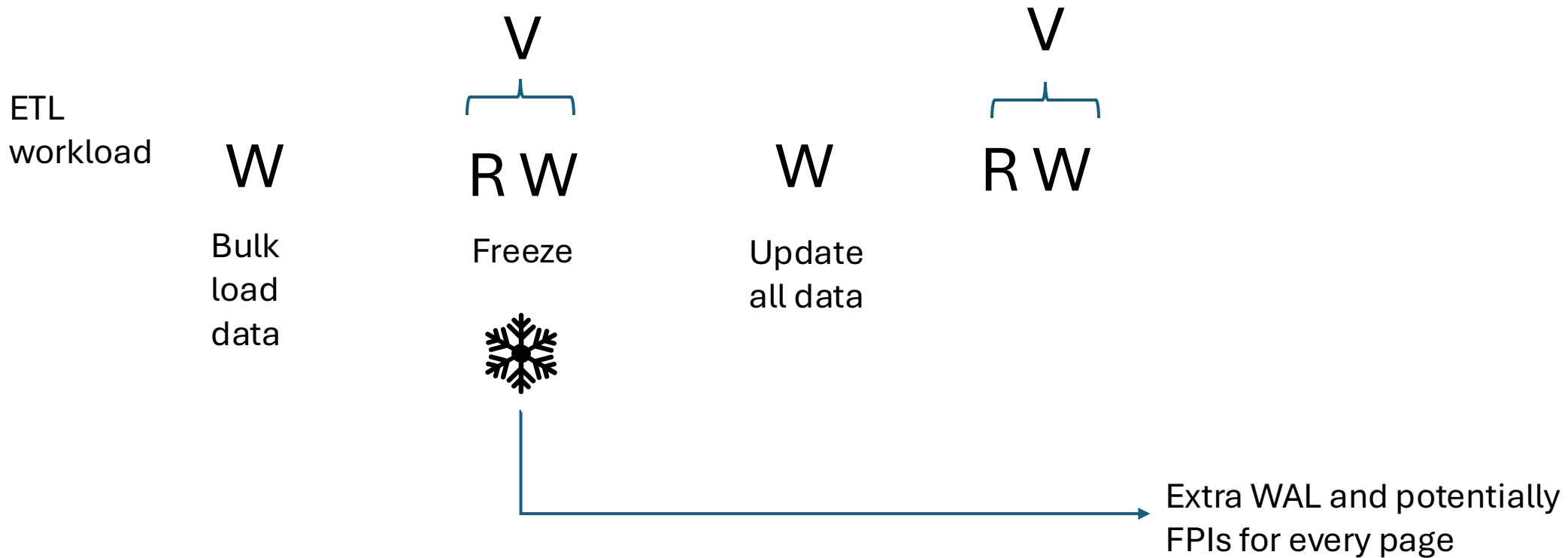
Hotly updated table



What about freezing every page once

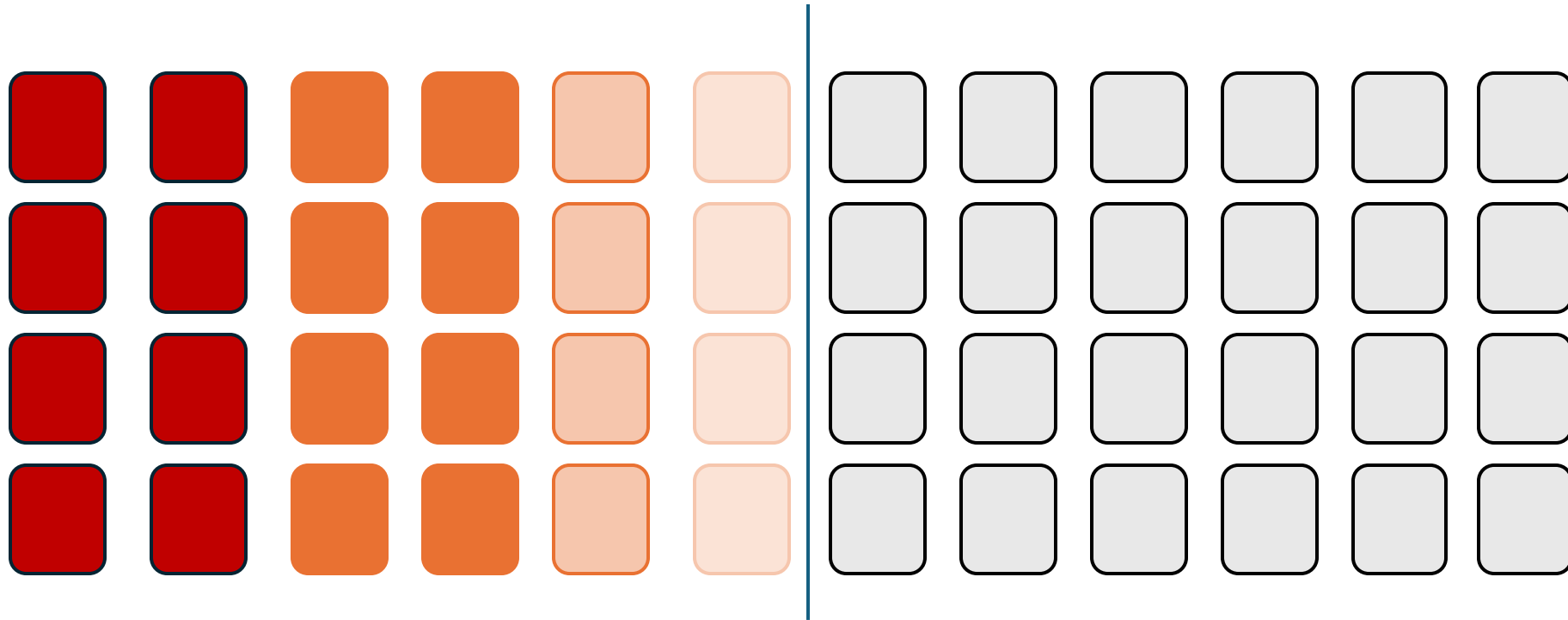


~~What about freezing every page once~~



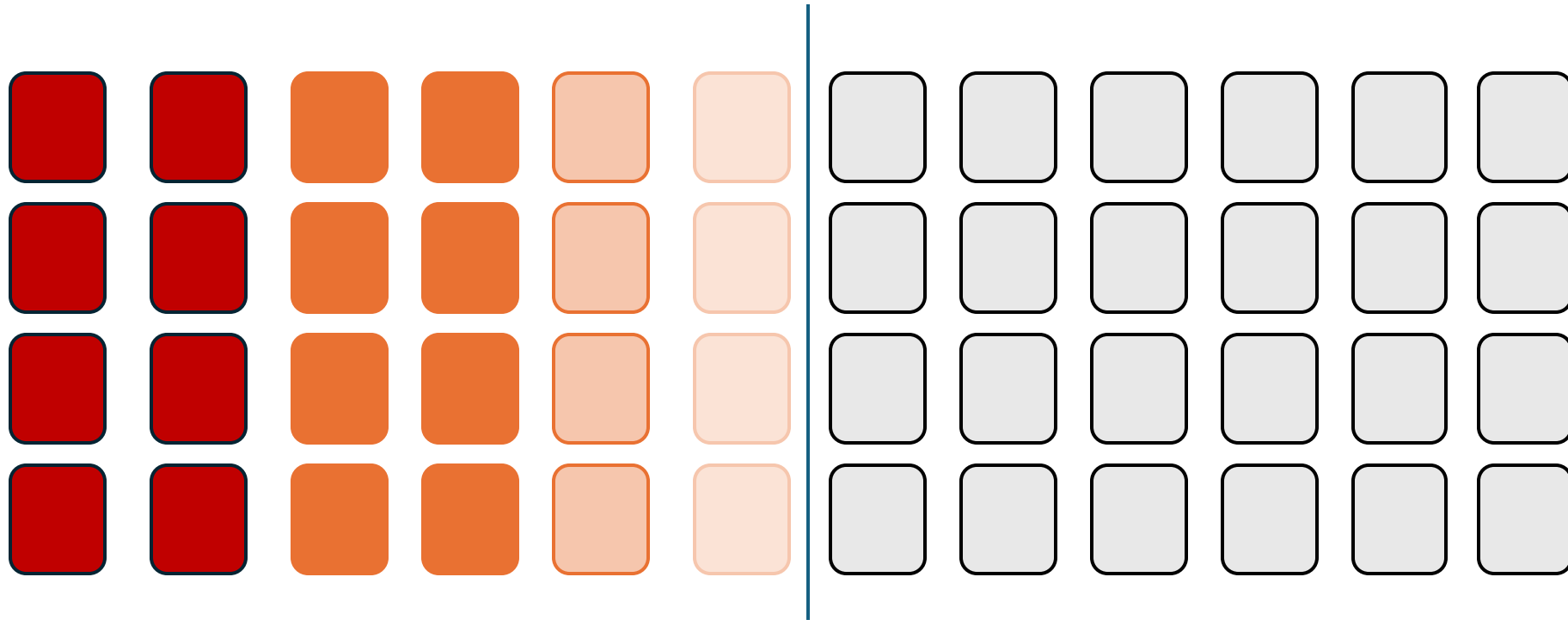
Need a model for table access
pattern and data age

Quiescence theory



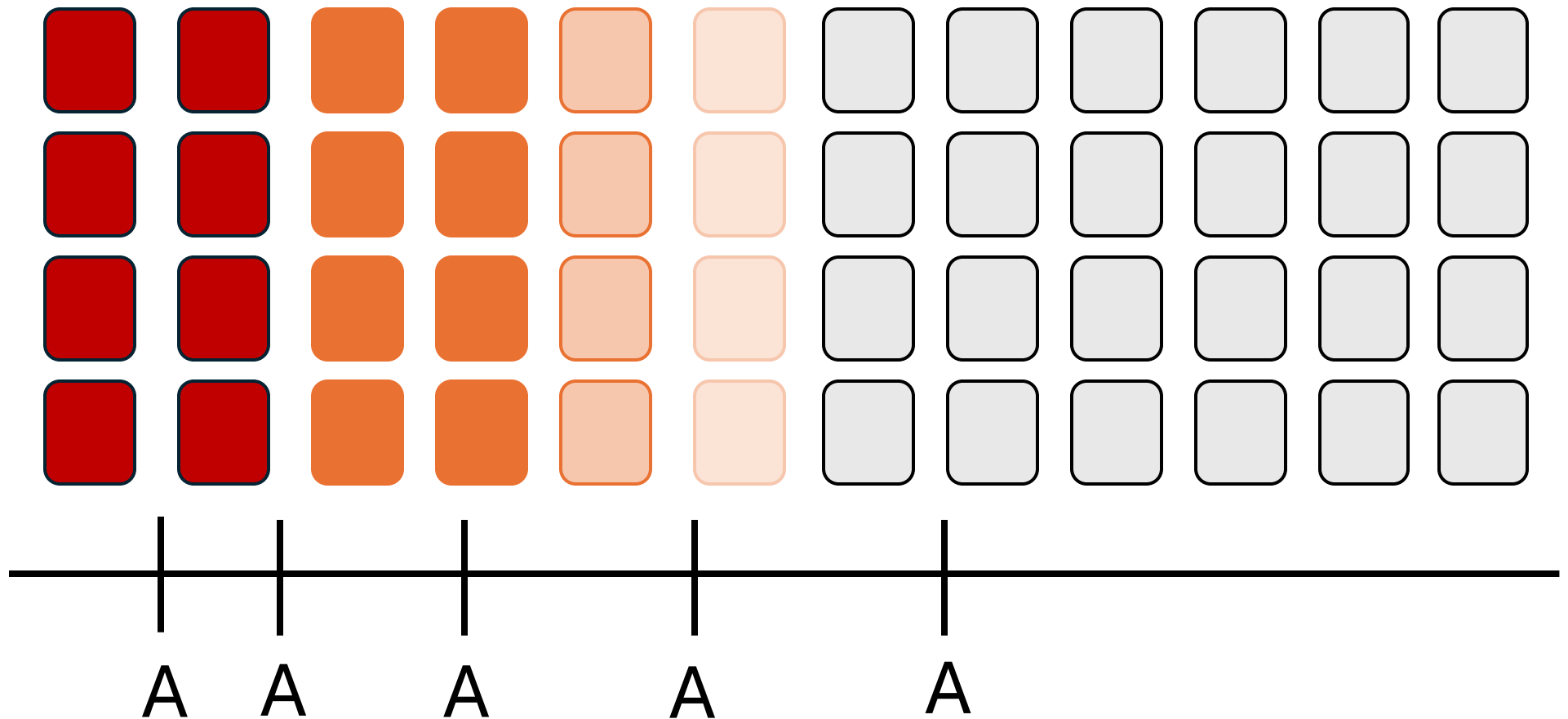
Evicted

Will the page be modified again?



Evicted

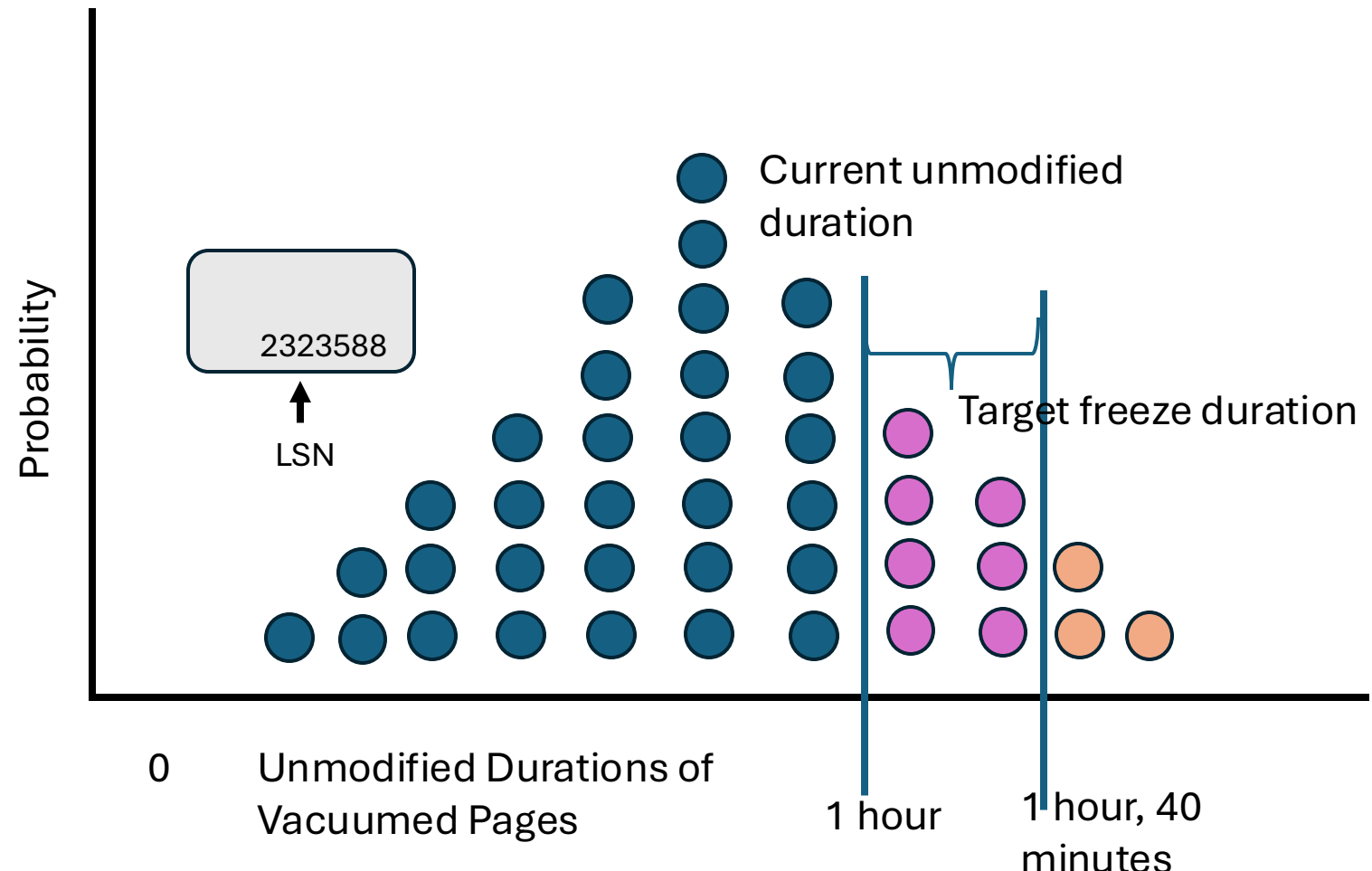
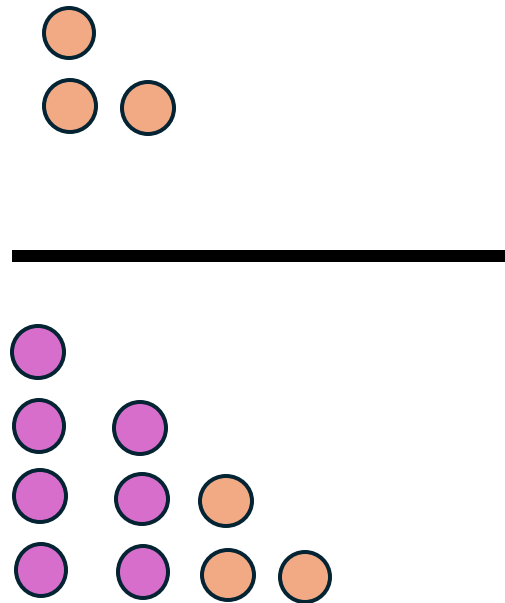
How often can you tolerate useless freezing?



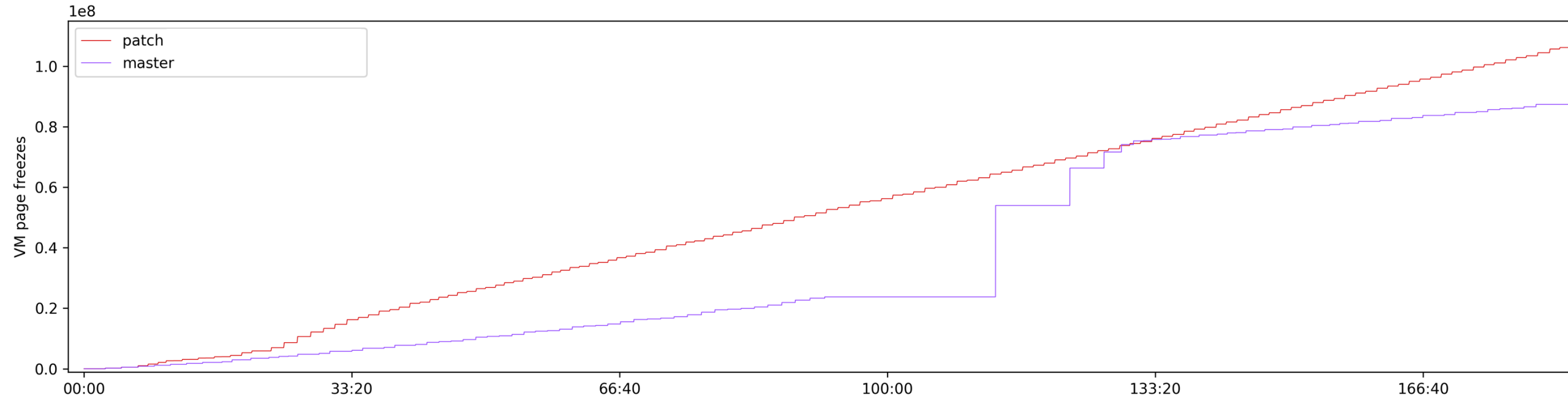
A = Autovacuum

Is the page young enough to stay
unmodified for target freeze
duration?

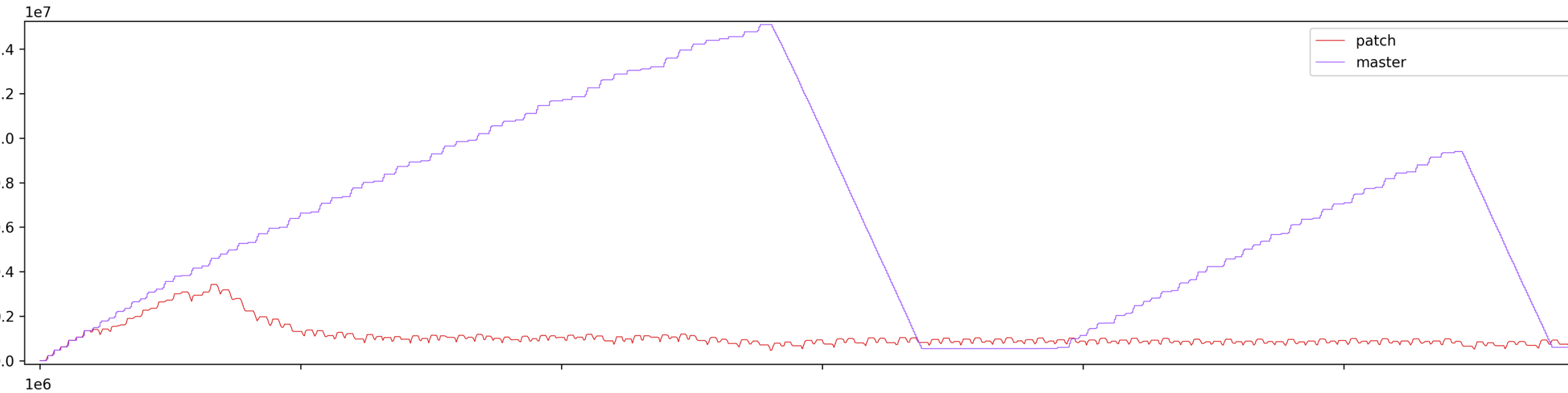
How likely is our specific page to stay unmodified for target freeze duration?



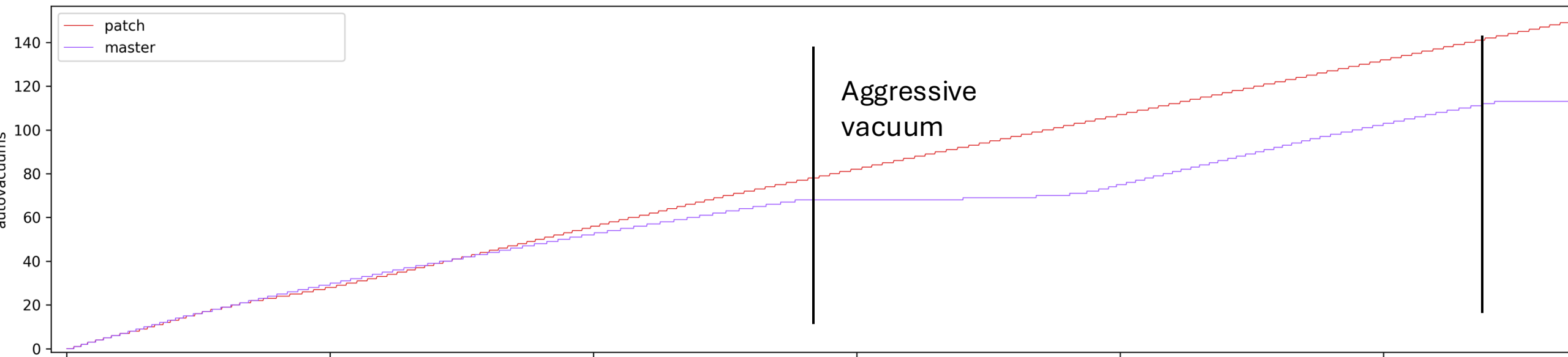
Results: Consistent freezing



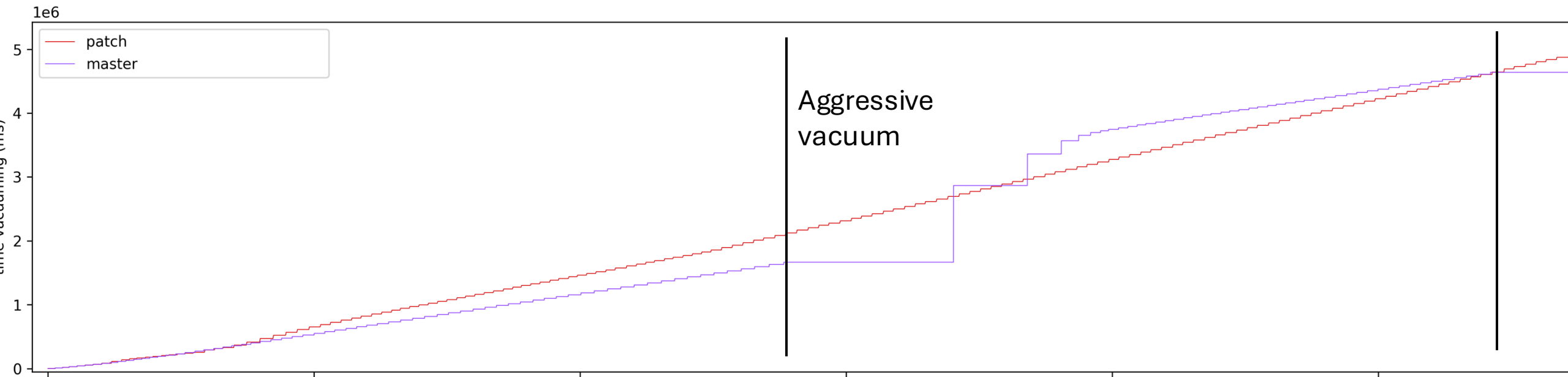
Results: All-visible Debt Low and Stable



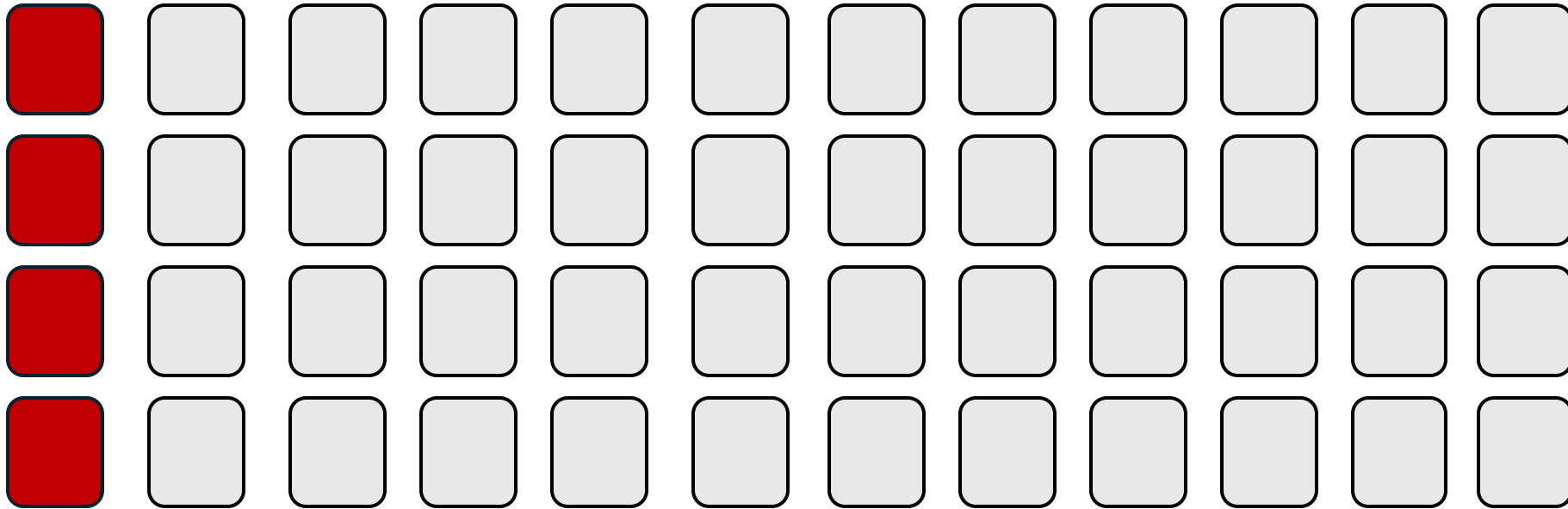
Many short autovacuum



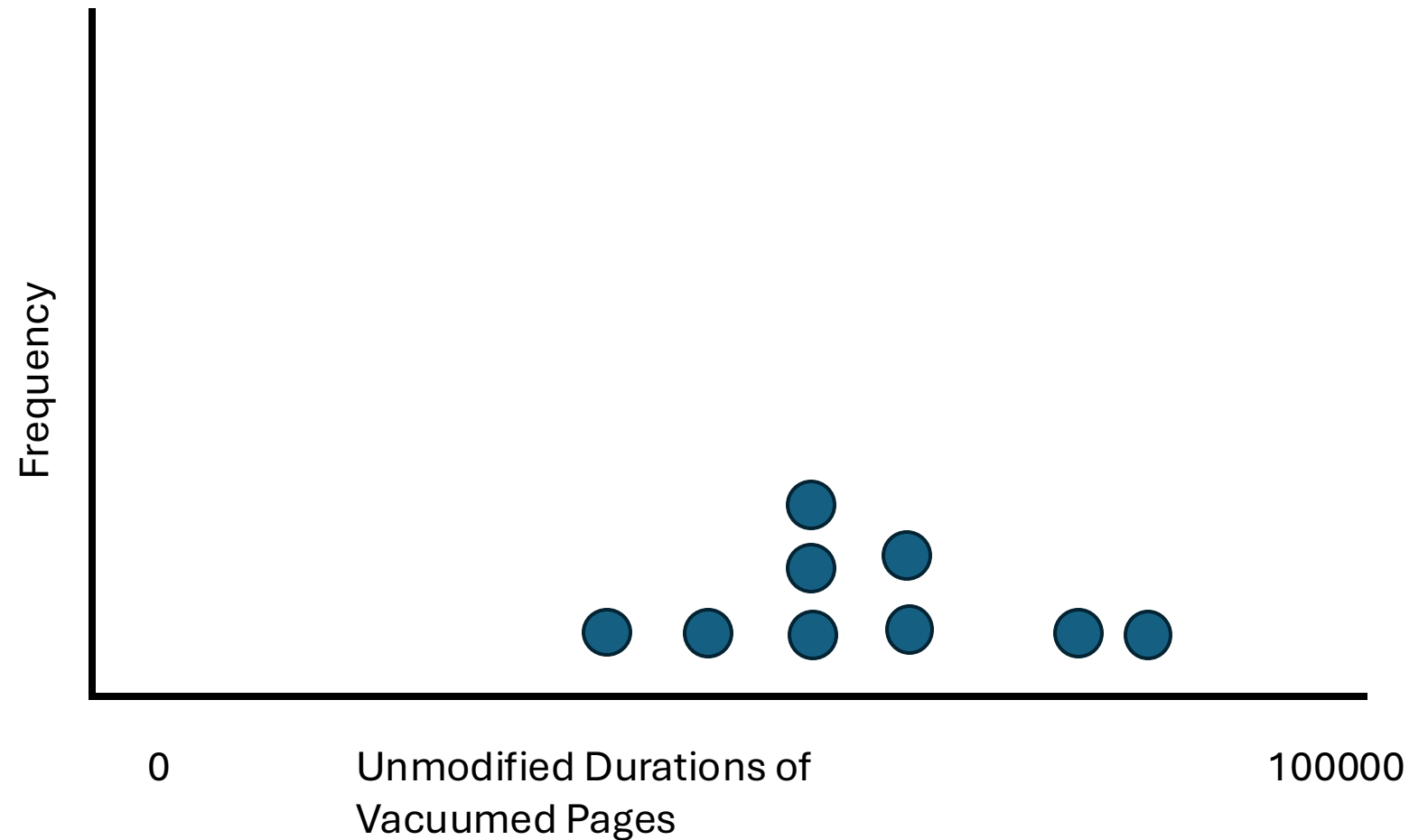
Lower Total Time Spent Vacuuming



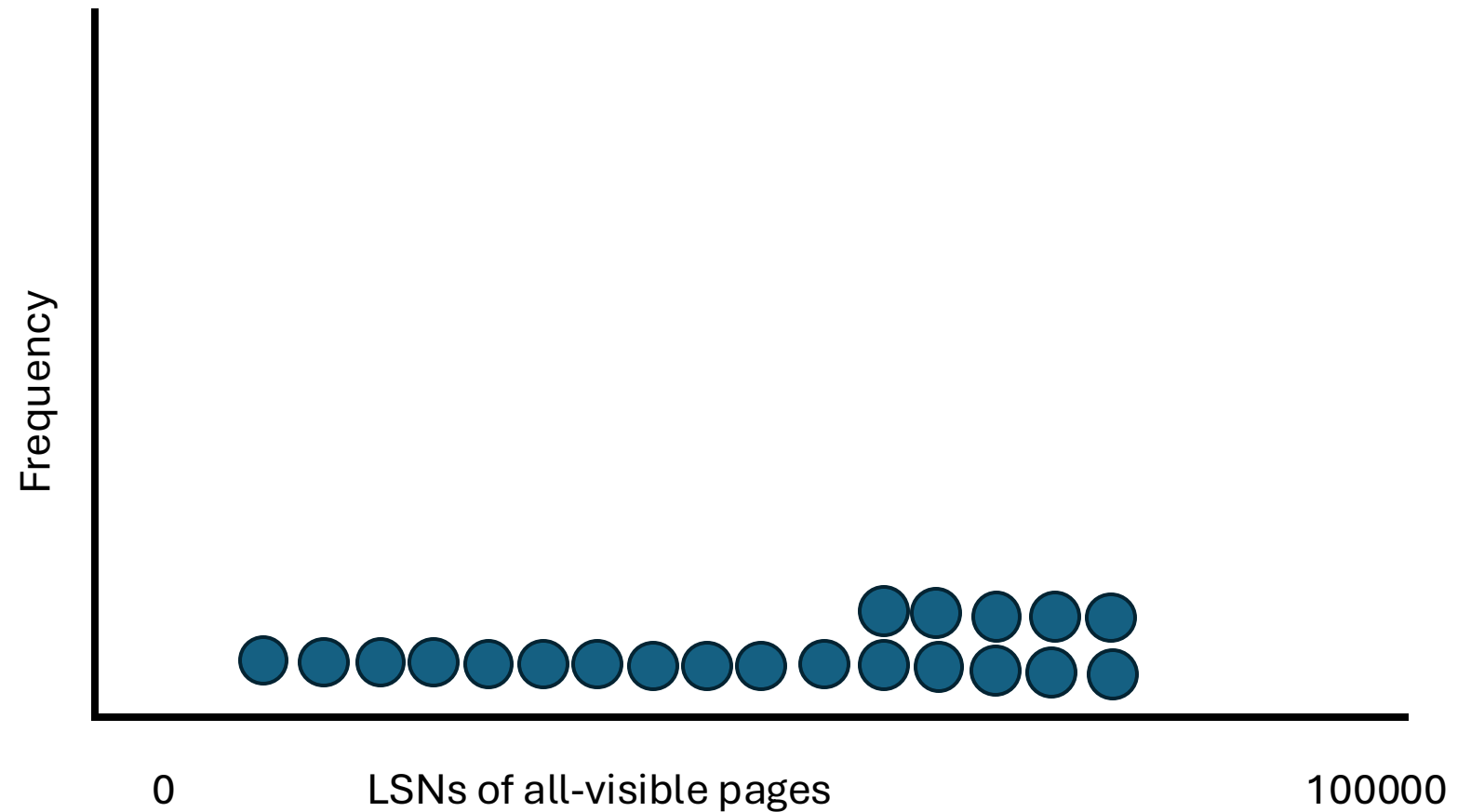
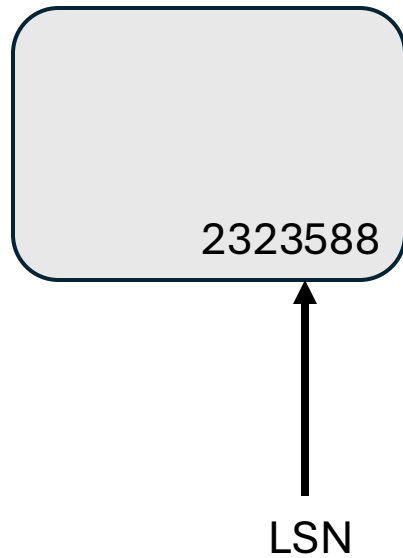
Insert-only workloads, data not modified



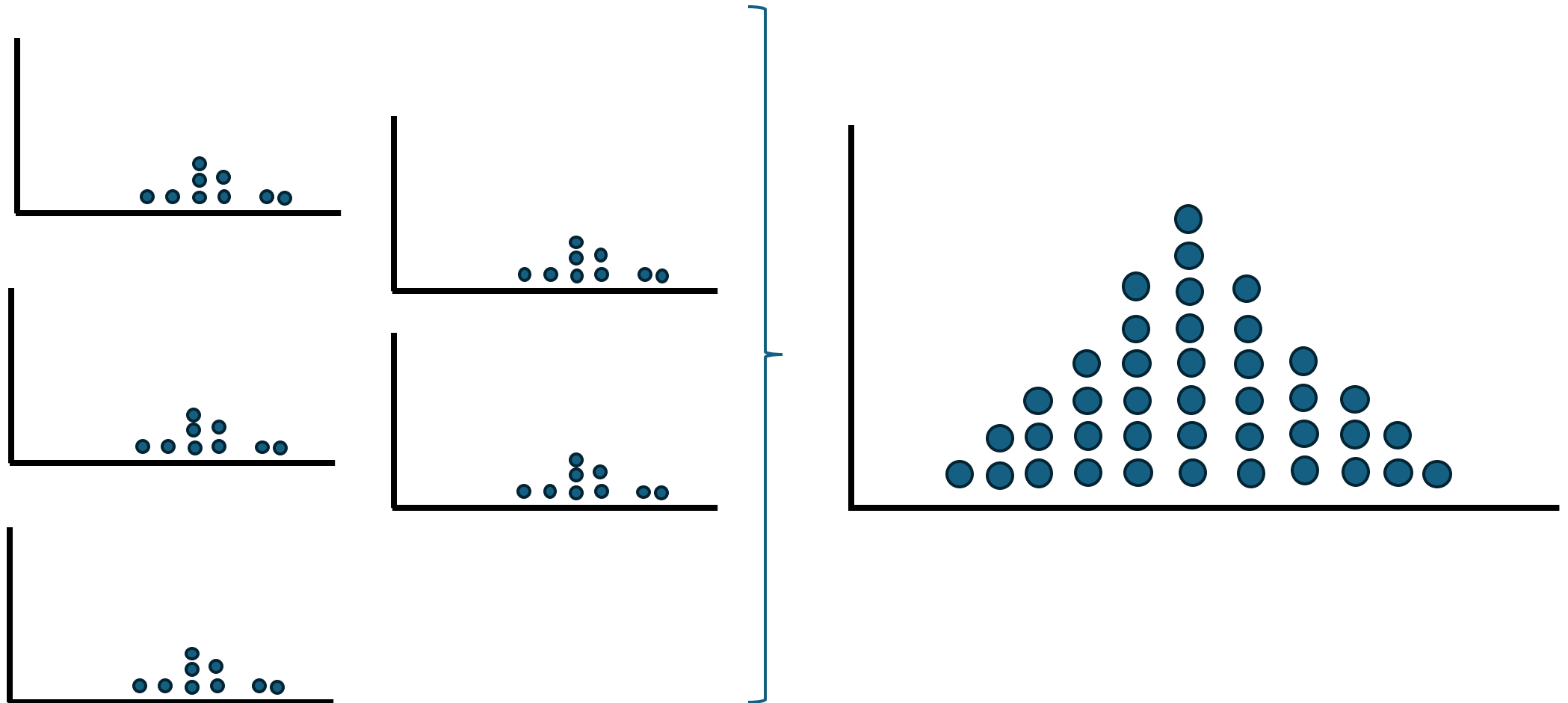
Missing data for pages not modified again



Need to add LSNs of all-visible pages



Building a distribution was complicated



High code complexity

Lots of code

- Low reusability

Major existing infrastructure issues

- Didn't work with failover or crash

Lessons Learned

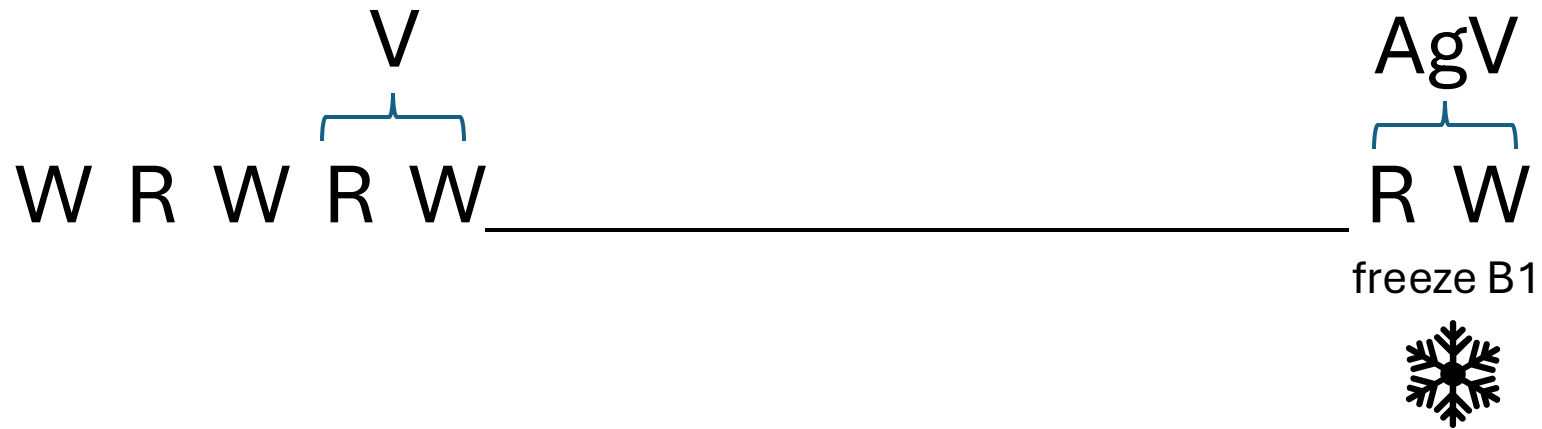


- Sometimes attempts to simplify fail
 - Define the problem better sooner
- 

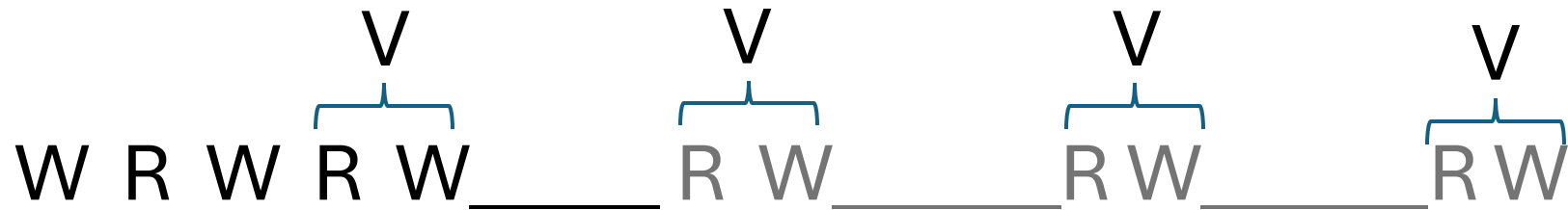
New Direction

What Went into Postgres 18

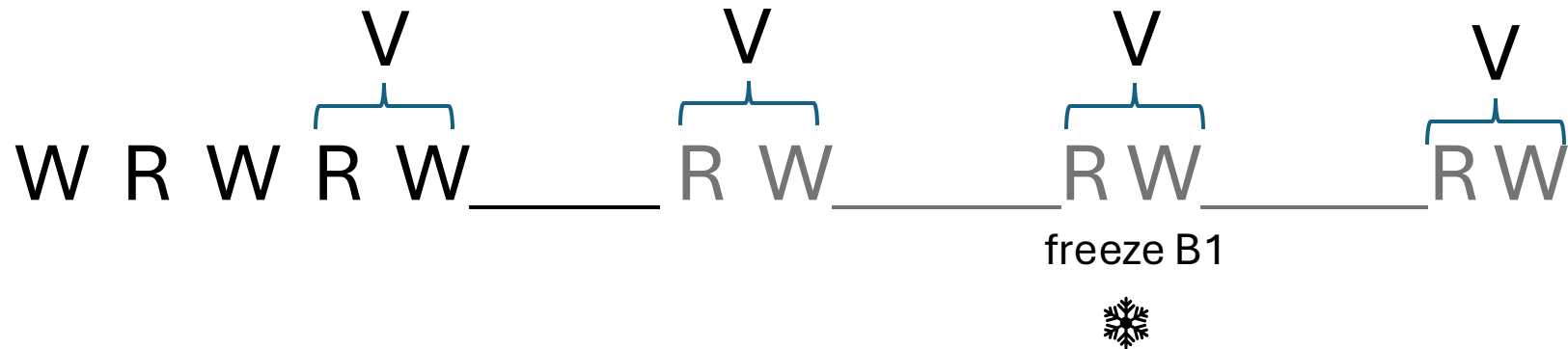
Long time between vacuuming and aggressive vacuuming B1



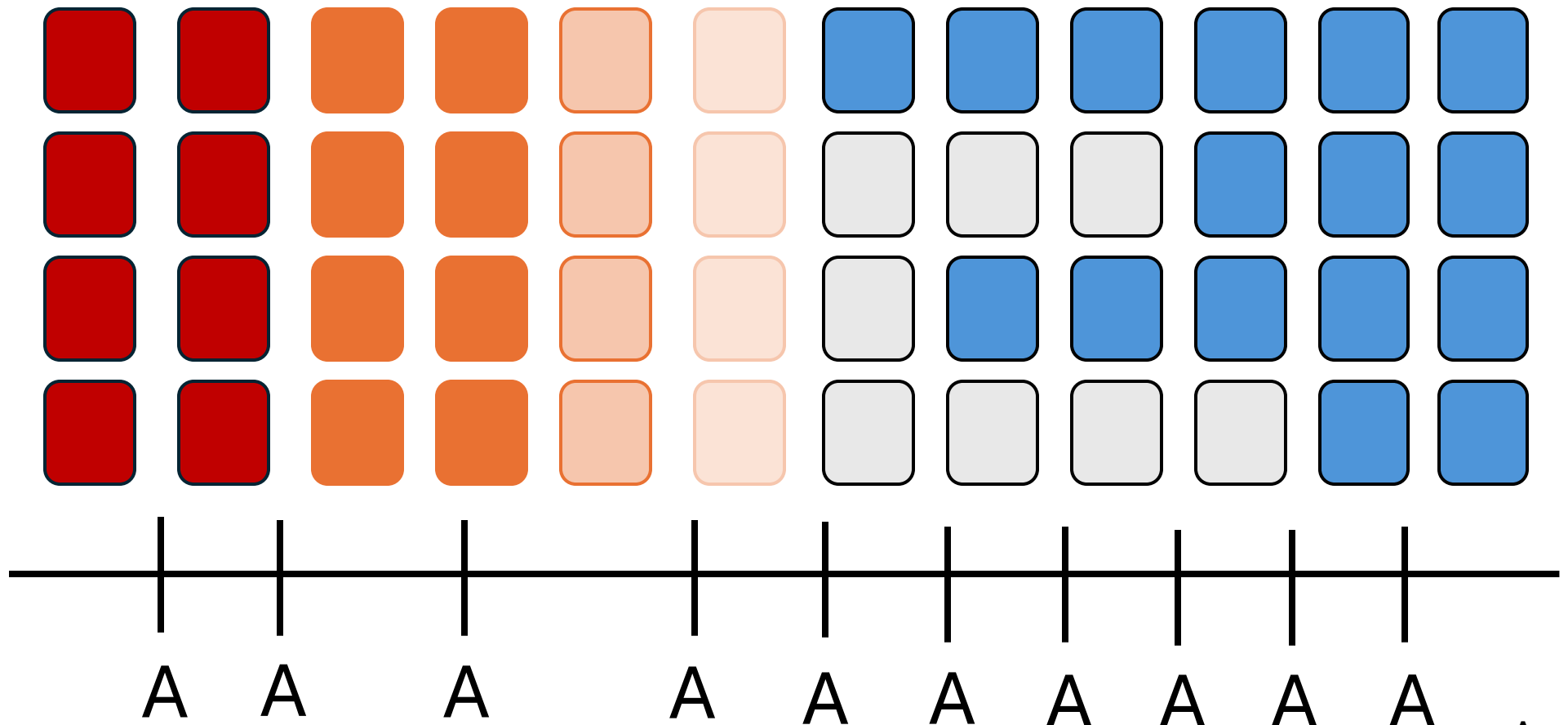
Other inserts triggered normal vacuums of new pages



What if we scan and freeze B1 sooner?



Switch framing to all-visible pages scanning



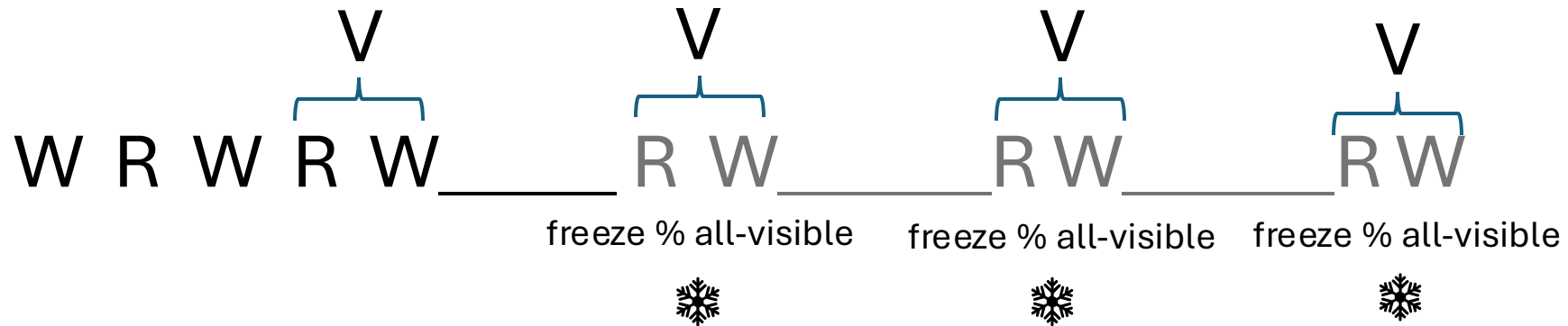
A = Autovacuum



Adaptively eager scan all-visible pages

- All-visible pages more likely to need freezing
 - Stop freezing if it isn't working
 - Only requires tracking information in one vacuum
-

Amortize the aggressive vacuum



Adaptively eager scan all-visible pages

- Cap eager scanning at a percentage of all-visible pages
- Suspend eager scanning if not successfully freezing pages



GUC and table storage option

- Eagerly scan and freeze up to 20% of the table
 - No more eager scanning for remainder of vacuum
 - `vacuum_max_eager_freeze_failure_rate`
 - Eagerly scan and fail to freeze 3% of 4096 block size region (32 MB)
 - Suspend eager scanning until next region
-

Future Directions

Evict inserted
data



Write

SELECT * queries
and sets hints



Read

Set
hints



C

Freeze on flush in checkpointer or bgwriter

Freeze on-
access

Evict hinted
data



Write

Normal
vacuum



Read

Set page vis
hint and VM



PD_ALL_VISIBLE



Vacuum
strategy evict



Write

Aggressive
vacuum



Read

Freeze tuples
and update VM



Vacuum
strategy evict

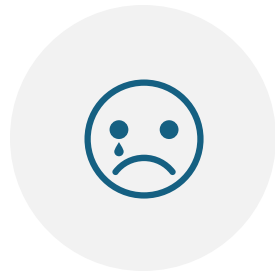


Write

Conclusion



ADAPTIVE ALGORITHMS
ARE POSSIBLE IN
POSTGRES



BUT ARE VERY HARD



BUT WE SHOULD THINK
MORE ABOUT THEM



THANKS TO
ANDRES FREUND AND
ROBERT HAAS



Got 3 minutes?
We'd love your input
on some of our
Postgres work



Get your FREE socks @ Microsoft booth